

# Information Theory II: mutual information and efficient coding

NEU 314, Fall 2021  
Lecture 18



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# Entropy

$$H(x) = - \sum_x p(x) \log p(x) \quad \text{in "bits"}$$
$$= \mathbb{E}[-\log p(x)]$$

other ways of writing it:

$$H(x) = \sum_{i=1}^N p_i (-\log p_i)$$
$$= \sum_{i=1}^N p_i \log \left( \frac{1}{p_i} \right)$$

how often it's used

code length / # questions

- average number of “yes/no” questions needed to identify  $x$
- average “surprise” from encountering a sample from  $p(x)$

# Conditional Entropy

$$H(x|y) = - \sum_y p(y) \sum_x p(x|y) \log p(x|y)$$

averaged  
over  $p(y)$

entropy of  $x$  given  
some fixed value of  $y$

# Conditional Entropy

$$H(x|y) = - \underbrace{\sum_y p(y)}_{\text{averaged over } p(y)} \underbrace{\sum_x p(x|y) \log p(x|y)}_{\text{entropy of } x \text{ given some fixed value of } y}$$

$$= - \sum_{x,y} p(x,y) \log p(x|y)$$

“On average, how uncertain are you about  $x$  if you know  $y$ ?”

“On average, how many questions do you need to identify  $x$  when you know  $y$ ?”

# exercise

Compute the conditional entropy:

X	1	2	3	4	5	6	7	8
$p(X Y=0)$	1/4	0	0	1/2	1/4	0	0	0
$p(X Y=1)$	0	0	0	0	0	0	1	0
Y	0	1						
$p(Y)$	2/3	1/3						

$$H(p(X|Y=0)) = 3/2$$

$$H(p(X|Y=1)) = 0$$

$$H(X | Y) = 2/3 (3/2) + 1/3 (0) = 1 \text{ bit}$$

“On average, you need 1 question to guess X when you know Y”

# Mutual Information

$$I(x, y) = H(x) - H(x|y)$$

total entropy in X minus  
conditional entropy of X given Y

$$= H(y) - H(y|x)$$

total entropy in Y minus  
conditional entropy of Y given X

$$= H(x) + H(y) - H(x, y)$$

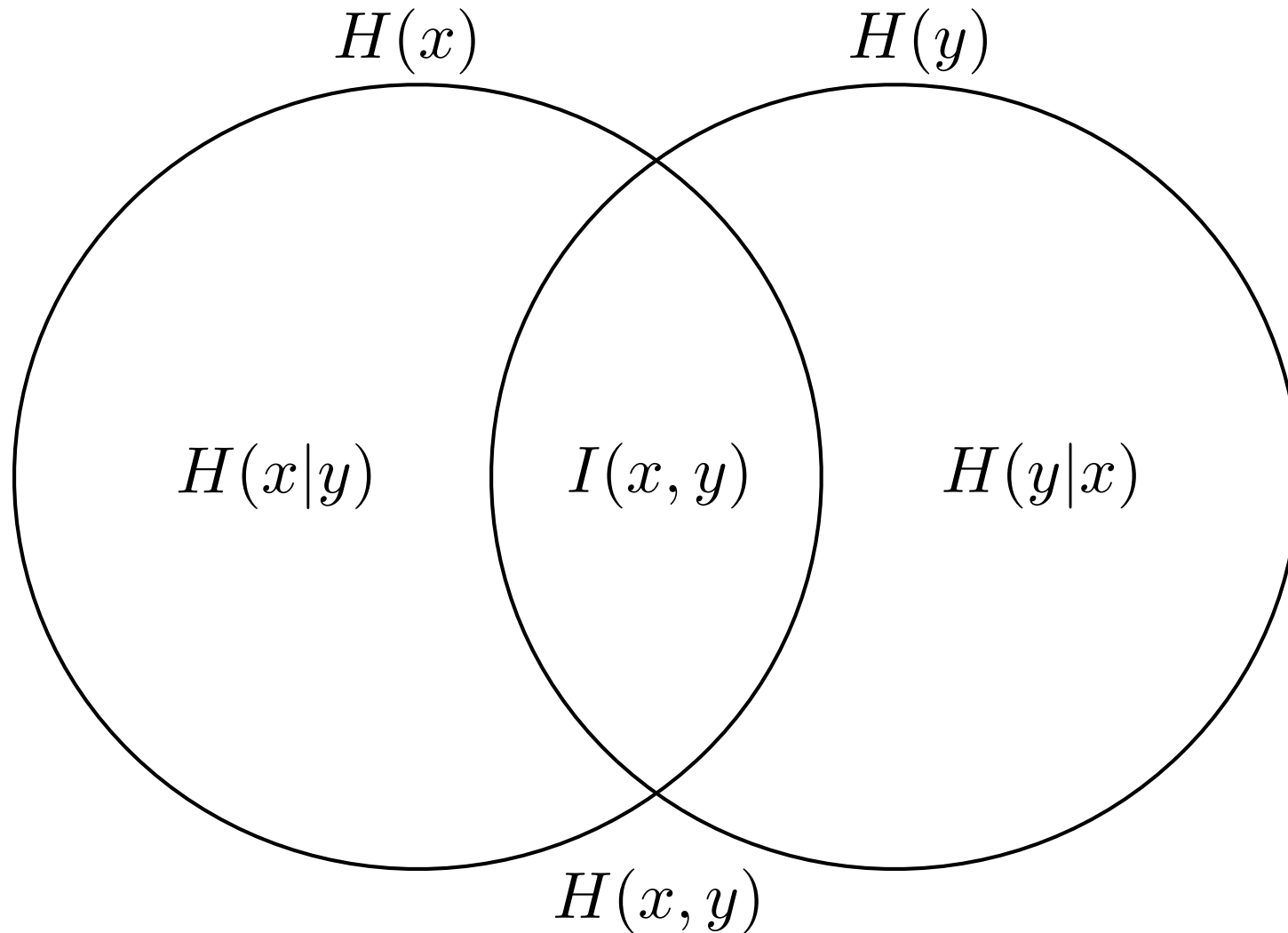
sum of entropies  
minus joint entropy

“How much does X tell me about Y (or vice versa)?”

“How much is your uncertainty about X reduced from knowing Y?”

“What is the difference between (# of questions needed to guess X) and (# questions needed to guess X when you’re given Y)”

# Venn diagram of entropy and information



# Kullback-Leibler Divergence

for two distributions  $P(x)$  and  $Q(x)$

$$D_{KL}(P||Q) = \sum_x P(x) \log \frac{P(x)}{Q(x)}$$
$$= \sum P(x) \log P(x) - \sum P(x) \log Q(x)$$
$$= \underbrace{\sum P(x)(-\log Q(x))}_{\substack{\text{avg under} \\ P(x)}} - \underbrace{\sum P(x)(-\log P(x))}_{\text{entropy of } P(x)}$$

code length based on  $Q(x)$       “cross-entropy”

• quantifies the number of *extra* bits required to code samples from  $P(x)$  if you use a codebook (“question asking strategy”) based on  $Q(x)$

Properties:

- $D_{KL}(P||Q) \geq 0, \forall P, Q$
- $D_{KL}(P||Q) = 0$ , iff  $P = Q$
- KL is not in general symmetric:  $D_{KL}(P||Q) \neq D_{KL}(Q||P)$



# Illustrating non-symmetry of KL divergence

$$D_{KL}(P||Q) = \sum_x P(x) \log \frac{P(x)}{Q(x)}$$

1<sup>st</sup> probability distribution  $P_1(X)$

1/8	1/8	1/8	1/8	1/8	1/8	1/8	1/8	1/8
1	2	3	4	5	6	7	8	

2<sup>nd</sup> probability distribution:  $P_2(X)$

0	0	0	1/8	1/8	1/8	1/8	1/2	
1	2	3	4	5	6	7	8	

**Exercise:** 1) What is  $KL(P_2 || P_1)$ ?

2) What is  $KL(P_1 || P_2)$ ?

# Mutual Information identities

$$\begin{aligned} I(x, y) &= H(x) - H(x|y) \\ &= - \sum_x p(x) \log p(x) + \sum_{x,y} p(x, y) \log p(x|y) \\ &= - \sum_{x,y} p(x, y) \log p(x) + \sum_{x,y} p(x, y) \log p(x|y) \\ &= \sum_{x,y} p(x, y) \log \frac{p(x|y)}{p(x)} \\ &= \sum_{x,y} p(x, y) \log \frac{p(x, y)}{p(x)p(y)} = D_{KL}(p(x, y) \| p(x)p(y)) \\ &= \sum_{x,y} p(x, y) \log \frac{p(y|x)}{p(y)} \quad \text{KL divergence between joint distribution} \\ &\quad \text{and product of marginals} \\ &= - \sum_{x,y} p(x, y) \log p(y) + \sum_{x,y} p(x, y) \log p(y|x) \\ &= - \sum_y p(y) \log p(y) + \sum_{x,y} p(x, y) \log p(y|x) \\ &= H(y) - H(y|x) \end{aligned}$$

# Data Processing Inequality

Suppose  $S \rightarrow R_1 \rightarrow R_2$  form a Markov chain, that is

$$P(R_1, R_2|S) = P(R_2|R_1)P(R_1|S)$$

Then necessarily:  $I(S, R_2) \leq I(S, R_1)$

- in other words, we can only lose information during processing

# Summary with formulas:

“surprise” function:  $-\log[p(x)]$

Entropy: “avg # Y/N Q’s” =  $-\sum_x P(x) \log P(x)$  (in bits)

Conditional Entropy:  $H(x|y) = -\sum P(x, y) \log P(x|y)$

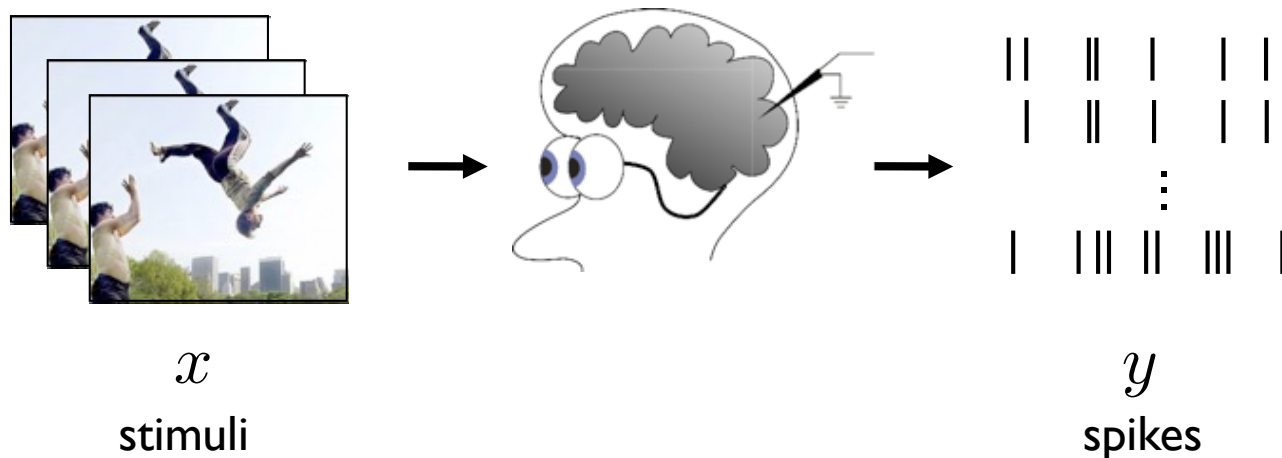
Mutual information:  $I(x, y) = H(x) - H(x|y) = H(y) - H(y|x)$   
 $= KL[p(x, y) \| p(x)p(y)]$

KL divergence  $D_{KL}(P||Q) = \sum_x P(x) \log \frac{P(x)}{Q(x)}$

Barlow's  
"Efficient Coding  
Hypothesis"

# Efficient Coding Hypothesis:

- goal of nervous system: maximize information about environment  
(one of the core “big ideas” in theoretical neuroscience)



## mutual information:

$$I(x, y) = \underbrace{H(y)}_{\text{response entropy}} - \underbrace{H(y|x)}_{\text{“noise” entropy}}$$

- avg # yes/no questions you can answer about  $x$  given  $y$  (“bits”)

# Barlow's original version:

## mutual information:

$$I(x, y) = H(y) - \cancel{H(y|x)}$$

response entropy    "noise" entropy

if responses are noiseless

# Barlow's original version:

## mutual information:

$$I(x, y) = H(y) - \cancel{H(y|x)} \quad \text{noiseless system}$$

response entropy    "noise" entropy

⇒ brain should maximize response entropy

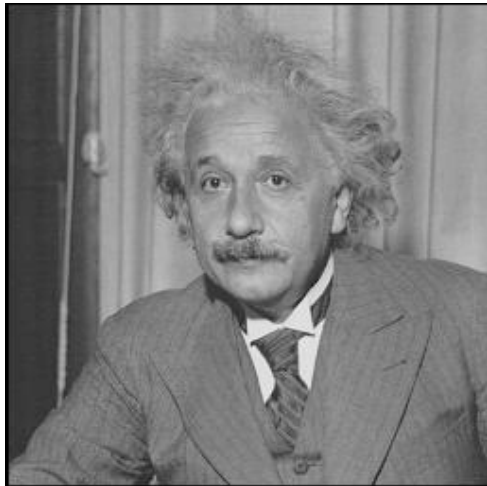
- use full dynamic range
- decorrelate ("reduce redundancy")

- mega impact: huge number of theory and experimental papers focused on decorrelation / information-maximizing codes in the brain



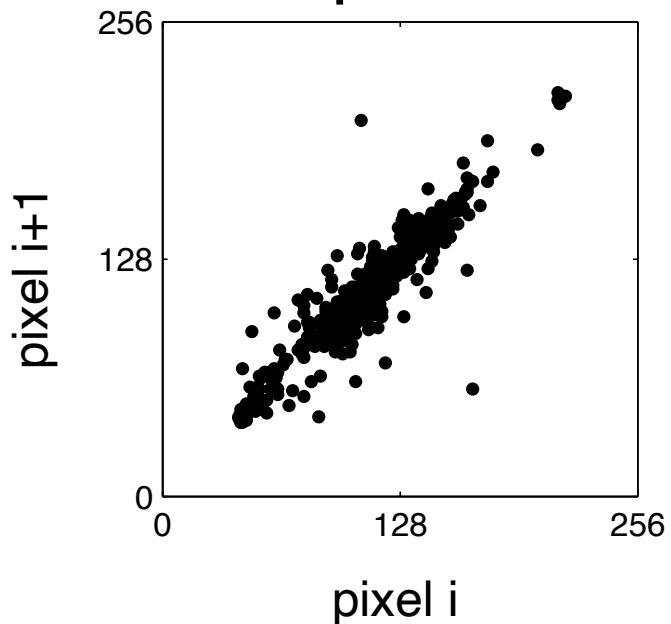
# basic intuition

natural image

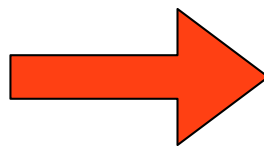


nearby pixels exhibit strong dependencies

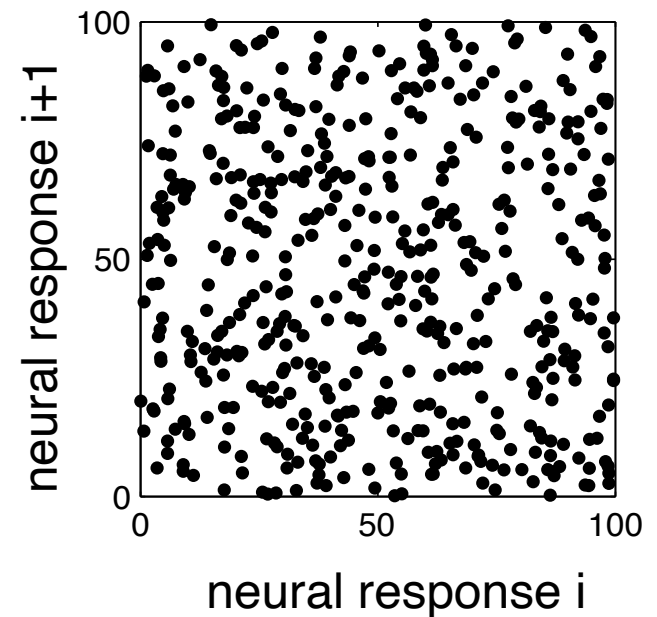
pixels



desired encoding



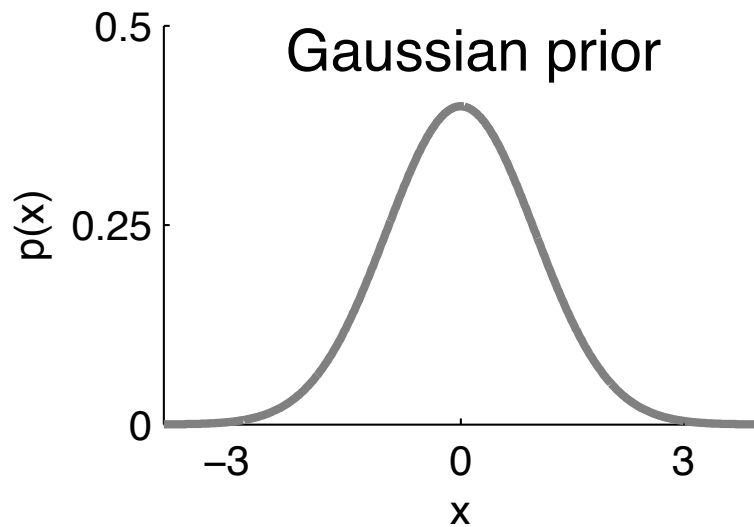
neural representation



# Application Example: single neuron encoding stimuli from a distribution $P(x)$

stimulus prior  $x \sim P(x)$

noiseless, discrete encoding  $y = f(x), \quad y \in \{y_1, y_2, \dots, y_n\}$

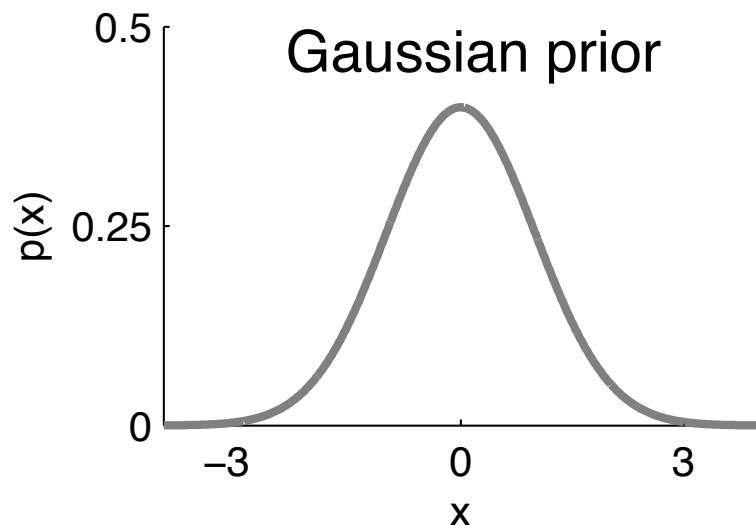


Q: what solution for infomax?

# Application Example: single neuron encoding stimuli from a distribution $P(x)$

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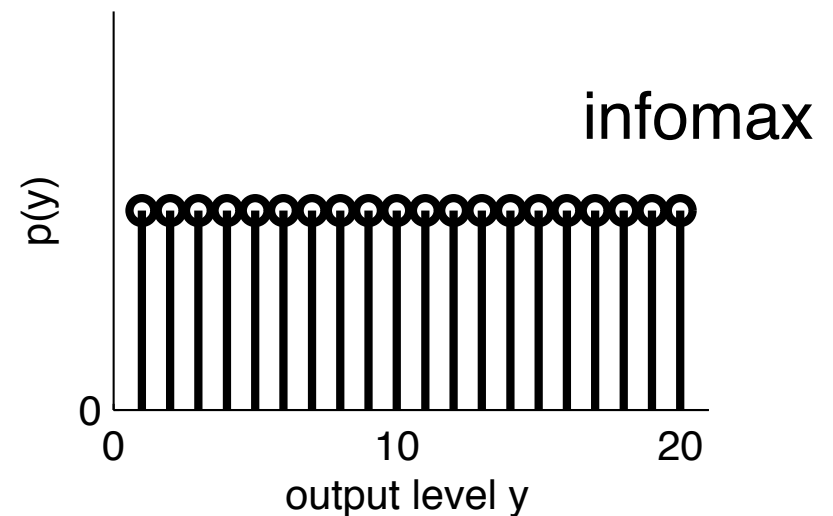
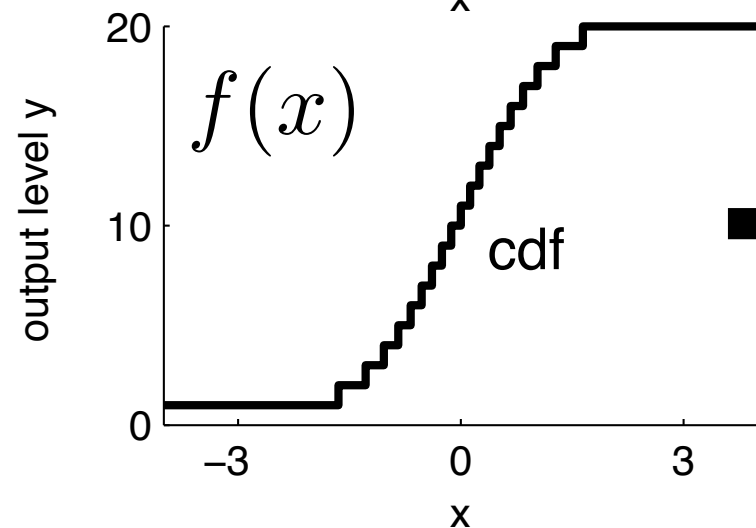
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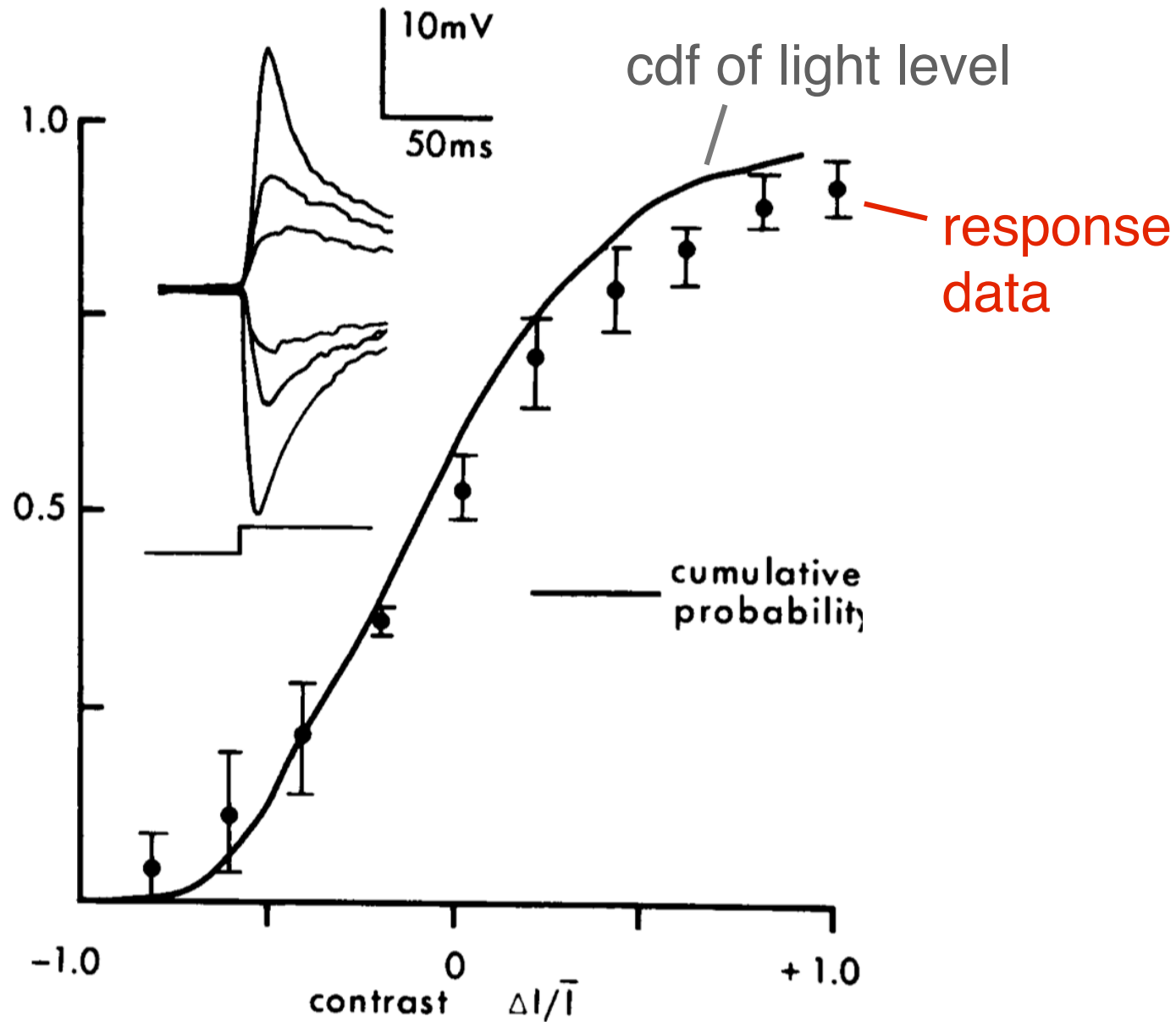
A: histogram-equalization

$$I(X, Y) = H(Y) - H(Y|X)$$



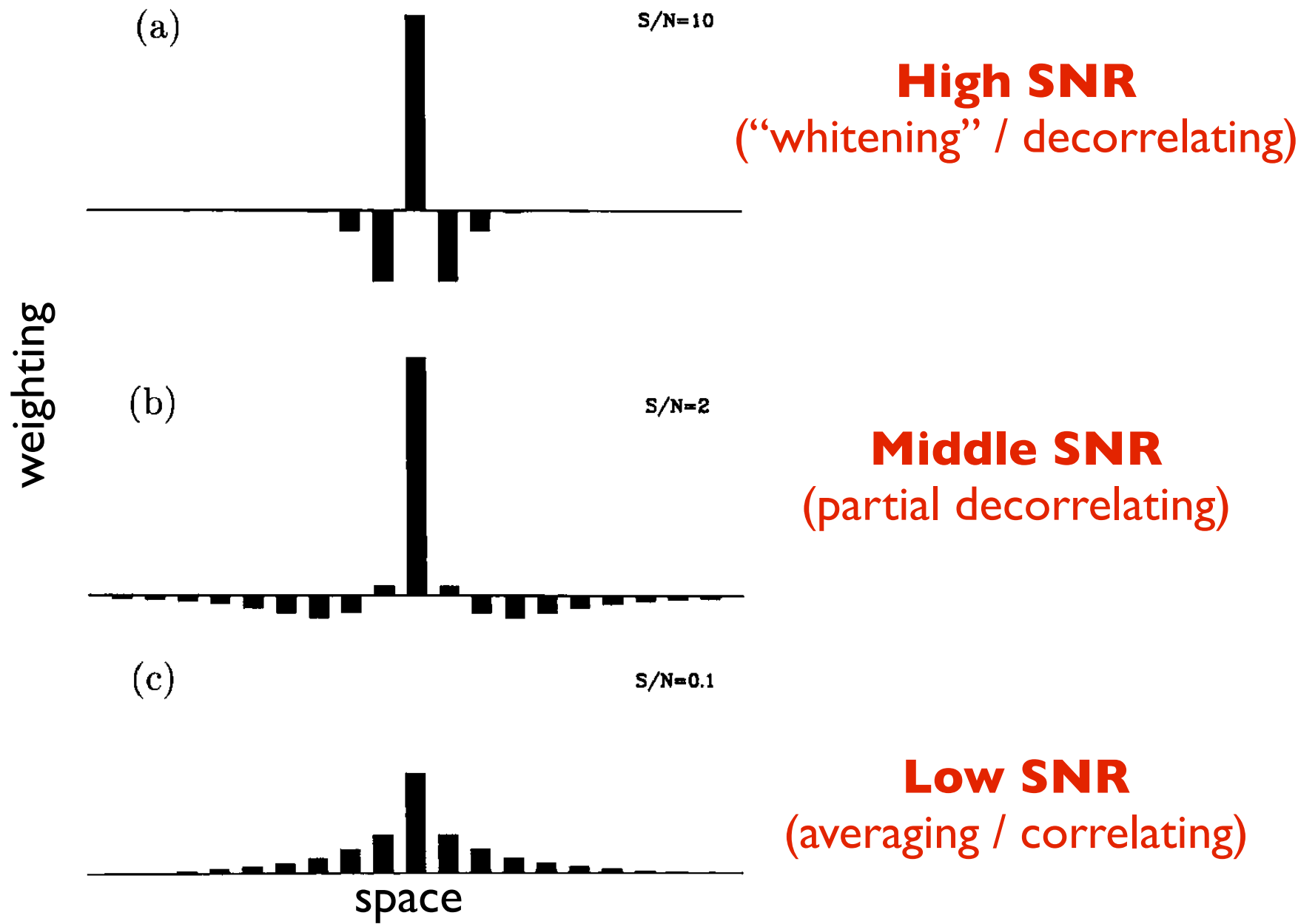
# Laughlin 1981: blowfly light response

- first major validation of Barlow's theory



# Atick & Redlich 1990 - extended theory to noisy responses

## luminance-dependent receptive fields



# summary: info theory

- entropy
- conditional entropy
- mutual information
- data processing inequality
- efficient coding hypothesis (Barlow)
  - neurons should “maximize their dynamic range”
  - multiple neurons: responses should decorrelate
  - Atick & Redlich: extended to noisy responses